Eliminating Disjunctions in Answer Set Programming by Restricted Unfolding

Jianmin Ji\(^1\), Hai Wan\(^2*\), Kewen Wang\(^3\), Zhe Wang\(^3\), Chuhan Zhang\(^2\), and Jiangtao Xu\(^2\)

\(^1\)School of Computer Science and Technology, University of Science and Technology of China, Hefei, China
\(^2\)School of Data and Computer Science, Sun Yat-sen University, Guangzhou, China
\(^3\)School of Information and Communication Technology, Griffith University, Griffith, Australia

\(\text{jianmin@ustc.edu.cn} \quad \text{wanhai@mail.sysu.edu.cn} \quad \text{k.wang,zhe.wang}@griffith.edu.au\)

Abstract

A disjunctive logic program under the answer set semantics can be equivalently translated to a normal logic program by the shifting transformation, if the program is head-cycle-free. In this paper, we provide an answer-set-preserving rewriting of a general disjunctive program to a normal program by first applying the unfolding transformation on atoms that prevent the program from being head-cycle-free, then shifting the resulting program. Different from other transformations that eliminate disjunctions in answer set programming, the new rewriting is efficient for “almost” head-cycle-free programs, i.e., programs that have only a few atoms that prevent them to be head-cycle-free. Based on the new rewriting, we provide an anytime algorithm to compute answer sets of a disjunctive program by calling solvers for normal logic programs. The experiment shows that the algorithm is efficient for some disjunctive programs. We also extend the rewriting to non-ground answer set programs on finite structures.

1 Introduction

Disjunctive logic programs extend normal logic programs by permitting disjunctions to appear in rule heads, which increases the expressive power of logic programs under the answer set semantics [Baral, 2003] and brings computational penalty as well. In particular, deciding whether a disjunctive program has an answer set is \(\Sigma^p_1\)-complete [Eiter and Gottlob, 1995] while deciding whether a normal program has an answer set is NP-complete. In practice, significant performance difference can be observed after adding disjunctive rules into normal programs, even if only a very small number of disjunctive rules are involved. Meanwhile, it is argued that many real-life problems can be characterized by programs consisting of a large number of normal rules and a few disjunctive rules. Hence, an interesting question is whether such a disjunctive program can be efficiently rewritten to an equivalent normal program under the answer set semantics.

Ben-Eliyahu and Dechter [1994] identified a class of disjunctive programs, called “Head-Cycle-Free” (HCF), and showed that each HCF program can be converted into an equivalent normal program in polynomial time by shifting head atoms into the body. The class of HCF programs is further generalised to that of Head-Elementary-loop-Free (HEF) programs [Gebser et al., 2011], which can be transformed into normal programs in polynomial time by shifting too. One could identify another proper super class of HEF programs that are polynomial time convertible to normal programs but it would be getting harder to check if a given disjunctive program is in such a class [Ji et al., 2015].

On the other hand, it is well known that a disjunctive program can be equivalently transformed into a negative disjunctive program by a set of program transformations including the unfolding transformation [Brass and Dix, 1997; Eiter and Wang, 2008]. Since shifting preserves the answer set semantics for negative disjunctive programs, each disjunctive program can be equivalently transformed into a normal program [Zhou, 2014]. Given that the unfolding transformation is exponential in the worst case, an exponential blow-up may occur in the unfolding-based rewriting. One source for this blow-up is from that unfolding can be applied on all atoms in the given disjunctive program. However, it has not been explored how the efficiency of unfolding can be improved in the context of disjunctive answer set programming. Especially, it is unclear if the unfolding transformation can be applied on only some atoms that are necessary.

A natural idea is to apply unfolding on atoms of the program one by one until the resulting program is HCF. However, unfolding may introduce new head-cycles and increase the set of atoms to be unfolded. Hence, generally speaking, one may need to unfold at all the atoms to achieve an HCF program, which indeed results in a negative program as in [Zhou, 2014]. In this paper, we show that we can restrict unfolding to “culprit” atoms, i.e., atoms that prevent the initial program to be HCF, and ignore new atoms that prevent the resulting program to be HCF. We prove that the result of our restricted unfolding, although not necessarily HCF, can be equivalently transformed to a normal program via shifting.

This result is of both theoretical and practical interests—it not only shows that the complexity penalty is closely tied to these “culprit” atoms, but also suggests the practicality of our rewriting approach on disjunctive programs that are “almost”...
HCF, i.e., with only a small number of “culprit” atoms. We make the following contributions:

- We provide an answer-set-preserving rewriting from a disjunctive program to a normal program by only unfolding at “culprit” atoms and shifting the resulting program, which is efficient for “almost” HCF programs.
- If II’ is a normal program rewritten from a disjunctive program II by restricted unfolding and shifting, we show that the computation of answer sets for II’ can be improved. In particular, although the rewriting may introduce new loops to II’, we only need to consider loops constructed from loops of II for computing answer sets of II’. This result suggests that although our rewriting to a normal logic program may lead to an exponential blow-up, there is not necessarily a blow-up in loop formulas.
- As another application of the rewriting, we present an anytime algorithm to compute answer sets of a disjunctive program II by calling solvers for normal logic programs. In each iteration, the algorithm outputs the answer sets of the shifted program II’ obtained from II by unfolding at a set A of atoms. The algorithm will output more answer sets of II when a larger A is applied.
- At last, we extend the rewriting to non-ground answer set programs on finite structures. As a result, we are able to compute answer sets of some first-order disjunctive logic programs by a solver for first-order normal programs such as [Asuncion et al., 2012].

2 Preliminaries

2.1 Disjunctive Logic Programs

We consider fully finite logic programs based on a propositional language $\mathcal{L}$ until Section 7. A (disjunctive) logic program (DLP) is a finite set of (disjunctive) rules of the form

$$a_1 \lor \cdots \lor a_k \leftarrow a_{k+1}, \ldots, a_m, \neg a_{m+1}, \ldots, \neg a_n.$$  \hspace{1cm} (1)

where $n \geq m \geq k \geq 1$ and $a_1, \ldots, a_m$ are atoms. If $k = 1$, it is a normal rule; if $k = m$, it is a negative rule; if $m = n$, it is a positive rule. A normal logic program (NLP) is a finite set of normal rules and a negative logic program (resp. positive logic program) is a finite set of negative (resp. positive) rules.

We will also write a rule $r$ of form (1) as

$$\text{head}(r) \leftarrow \text{body}(r).$$

where $\text{head}(r)$ is $a_1 \lor \cdots \lor a_k$, $\text{body}(r) = \text{body}^+(r) \lor \text{body}^-(r)$, $\text{body}^+(r)$ is $a_{k+1} \lor \cdots \lor a_m$, and $\text{body}^-(r)$ is $\neg a_{m+1} \lor \cdots \lor \neg a_n$. We define $\text{head}(r)$, $\text{body}^+(r)$, $\text{body}^-(r)$ with their corresponding sets of atoms, and $\text{body}(r)$ the set $\{a_{k+1}, \ldots, a_m, \neg a_{m+1}, \ldots, \neg a_n\}$. Given a DLP II, we use $\text{Atoms}(\Pi)$ to denote the set of atoms occurring in II. We use $\text{tr}(r)$ to denote the propositional formula $\text{body}(r) \supset \text{head}(r)$ and $\text{tr}(\Pi) = \bigwedge_{r \in \Pi} \text{tr}(r)$.

A set $S$ of atoms satisfies a rule $r$ if $S \cup \{\neg p \mid p \notin S\} \models \text{tr}(r)$ in the sense of propositional logic. $S$ satisfies a program II if $S$ satisfies every rule in II.

The answer sets of a DLP are defined in [Gelfond and Lifschitz, 1991]. Given a DLP II and a set $S$ of atoms, the Gelfond-Lifschitz reduct of II on $S$, written $\Pi^S$, is obtained from II by deleting:

1. each rule that has a formula $\neg p$ in its body with $p \in S$,
2. all formulas of the form $\neg p$ in the bodies of the remaining rules.

A set $S$ of atoms is an answer set of II if $S$ is a minimal set satisfying $\Pi^S$. We use $\text{AS}(\Pi)$ to denote the set of answer sets of II.

2.2 Loops and Loop Formulas

Lee and Lifschitz [2003] extended the notions of loops and loop formulas [Lin and Zhao, 2004] to DLPs. For a DLP II, the positive dependency graph of II, written $G_{II}$, is the directed graph whose vertices are atoms in II, and there is an arc from $p$ to $q$ if there is a rule $r$ in II such that $p \in \text{head}(r)$ and $q \in \text{body}^+(r)$. A set $L$ of atoms is a loop of II if the $L$-induced subgraph of $G_{II}$ is strongly connected. Note that, every singleton whose atom occurs in II is also a loop of II. We use $\text{Loop}(\Pi)$ to denote the set of loops of II.

Given a DLP II and a loop $L$ of II, a rule $r \in \Pi$ is an external support of $L$ under II if $\text{head}(r) \cap L \neq \emptyset$ and $L \cap \text{body}^+(r) = \emptyset$. Let $R^-(L, \Pi)$ be the set of external support rules of $L$ under II. The (conjunctive) loop formula of $L$ under II, written $\text{LF}(L, \Pi)$, is the following implication:

$$\Phi(L, \Pi) = \bigwedge_{p \in L} \bigvee_{r \in R^-(L, \Pi)} \left(\text{body}(r) \land \bigwedge_{q \in \text{head}(r) \cap L} \neg q\right).$$


Theorem 1 (Theorem 1 in [Lee and Lifschitz, 2003])

For a DLP II and a set $S$ of atoms, $S$ is an answer set of II if and only if $S \cup \{\neg p \mid p \notin S\}$ is a model of $\Phi(L, \Pi)$.

2.3 Unfolding and Shifting

We first review the unfolding transformation [Gergatsoulis, 1997; Brass and Dix, 1997].

Definition 1 (Elementary unfolding) Let $r_1$, $r_2$ be two rules and there is an atom $a \in \text{head}(r_1) \cap \text{body}^+(r_2)$. The rule obtained by elementary unfolding $r_2$ using $r_1$ at $a$, denoted by $\text{unfold}(r_1, r_2, a)$, is the rule $r$ such that $\text{head}(r) = (\text{head}(r_1) \setminus \{a\}) \cup \text{head}(r_2)$ and $\text{body}(r) = \text{body}(r_1) \cup (\text{body}(r_2) \setminus \{a\})$.

Definition 2 (Unfolding rules at an atom) Let $\Pi$ be a DLP and $a$ an atom in $\Pi$. The set of rules obtained by unfolding rules at $a$ in II, denoted by $\text{unfold}_{\Pi}(a)$, is the set $\{\text{unfold}(r_1, r_2, a) \mid r_1, r_2 \in \Pi \land a \in \text{head}(r_1) \cap \text{body}^+(r_2)\}$.

Let $\Pi$ be a DLP and $a \in \text{Atoms}(\Pi)$. We denote

$$\text{remove}(\Pi, a) = \Pi \setminus \{r \mid r \in \Pi \land a \in \text{body}^+(r)\} \cup \text{unfold}_{\Pi}(a).$$

Example 1 Consider the logic program $\Pi_1$:

$$a \leftarrow b. \quad d \leftarrow b. \quad e \leftarrow a, d. \quad c \leftarrow e. \quad c \leftarrow a. \quad d \leftarrow b. \quad b \leftarrow a.$$
 unfold (a) is the set of rules: 
\[ \begin{align*}
\text{b} \lor \text{c} &\leftarrow \text{d}, \\
\text{e} &\leftarrow \text{b}, \\
\text{b} &\leftarrow \text{c}, \\
\text{b} &\leftarrow \text{b}.
\end{align*} \]

From Theorem 13 in Gergatsoulis, 1997, we have the following theorem.

**Theorem 2** Let \( \Pi \) be a DLP and \( a \) an atom in \( \Pi \). Then 
\[ A^S(\Pi) = A^S(\text{remove}(\Pi, a)) \]  

Given a DLP \( \Pi, E \subseteq \text{Atoms}(\Pi) \), and \( a \in E \), we can inductively define \( \text{remove}(\Pi, E) \) as follows:

- \( \text{remove}(\Pi, \{a\}) = \text{remove}(\Pi, a) \)
- \( \text{remove}(\Pi, E) = \text{remove}(\Pi, E \setminus \{a\}) \{a\} \)

**Corollary 3** Let \( \Pi \) be a DLP and \( E \subseteq \text{Atoms}(\Pi) \). Then 
\[ A^S(\Pi) = A^S(\text{remove}(\Pi, E)) \]  

Gelfond et al. [1991] provided a mapping from a DLP to an NLP by “shifting” head atoms into the bodies. We denote \( \text{sh}(\Pi) \) to be the NLP obtained from a DLP \( \Pi \) by substituting every rule of form \( (1) \) with the \( k \) rules

\[ \begin{align*}
a_1 &\leftarrow \neg a_1, \ldots, \neg a_{i-1}, \neg a_{i+1}, \ldots, \neg a_k, \\
a_k+1, \ldots, a_m, \neg a_{m+1}, \ldots, \neg a_n, \quad (1 \leq i \leq k)
\end{align*} \]

It is known that every answer set of \( \text{sh}(\Pi) \) is also an answer set of \( \Pi \), but the converse is not true in general. Ben-Eliyahu and Dechter [1994] identified a class of DLP, called “Head-Cycle-Free” (HCF), and showed that the converse is true if \( \Pi \) is an HCF program. In particular, a DLP \( \Pi \) is called HCF, if \( \text{head}(r) \cap L \leq 1 \) for every rule \( r \) of \( \Pi \) and every loop \( L \) of \( \Pi \), and the answer sets of an HCF program \( \Pi \) coincide with the answer sets of the NLP \( \text{sh}(\Pi) \).

Given a DLP \( \Pi \), \( \text{remove}(\Pi, \text{Atoms}(\Pi)) \) is equivalent to a negative program [Brass and Dix, 1997], which is HCF. Then 
\[ A^S(\Pi) = A^S(\text{remove}(\Pi, \text{Atoms}(\Pi))) = A^S(\text{sh}(\text{remove}(\Pi, \text{Atoms}(\Pi)))) \]  

Based on this result, Zhou [2014] provided a syntactic transformation from a DLP to an NLP by first applying the unfolding transformation on all the atoms, then shifting the resulting program to the NLP.

### 3 Eliminating Disjunctions by Restricted Unfolding

In this section, we show that a disjunctive program can be equivalently translated to a normal program by first applying the unfolding transformation only on “culprit” atoms, i.e., atoms that prevent the program from being HCF.

For a DLP \( \Pi \), we denote

\[ \text{HC}(\Pi) = \{ a \mid \text{there exist } r \in \Pi \text{ and } L \in \text{Loop}(\Pi) \text{ such that } |\text{head}(r) \cap L| > 1 \text{ and } a \in \text{head}(r) \cap L \} \]

Intuitively, \( \text{HC}(\Pi) \) specifies the set of atoms that prevent \( \Pi \) to be HCF.

Now we provide the lemmas that draw a sketch of the proof for our main result.

**Lemma 2** Let \( \Pi \) be a DLP and \( E \subseteq \text{Atoms}(\Pi) \), and \( L \subseteq \text{Loop}(\text{remove}(\Pi, E)) \). If \( L \cap \text{HC}(\Pi) = \emptyset \) then 
\[ \text{LF}(L, \text{remove}(\Pi, E)) = \text{LF}(L, \text{sh}(\text{remove}(\Pi, E))). \]

As discussed in Introduction, unfolding may introduce new head-cycles and hence increase the set of atoms in HC.

**Example 1 (Continued)** \{b, c, e\} is not a loop of \( \Pi_1 \), while it is a loop of the program \( \text{remove}(\Pi_1, \{a\}) \), then 
\[ \text{HC}(\text{remove}(\Pi_1, \{a\})) = \{b, e\} \text{ while } \text{HC}(\Pi_1) = \{a, b\} \]

For some DLP \( \Pi \) and \( E \subseteq \text{Atoms}(\Pi) \), exponentially many loops can be introduced by unfolding, and \( \text{HC}(\text{remove}(\Pi, E)) \) can be greatly larger than \( \text{HC}(\Pi) \).

**Example 2** Consider the logic program \( \Pi_2 \):

\[ a \lor b \leftarrow c, \quad e \leftarrow a, \quad a \leftarrow b. \]

\[ a_1 \lor \cdots \lor a_n \leftarrow b, \quad d \leftarrow a_i, \quad (1 \leq i \leq n) \]

where \( n \geq 1 \). Note that \( \quad \text{Loop}(\Pi_2) = \bigcup_{\{p \in \text{Atoms}(\Pi_2) \}} \{\{p\} \} \bigcup \{\{a, b\}\} \quad \text{HC}(\Pi_2) = \{a, b\} \quad \text{remove}(\Pi_2, a) = (\Pi_2 \setminus \{e \leftarrow a, d \leftarrow b \}) \bigcup \{b \lor e \leftarrow b \} \bigcup \{e \leftarrow d, \quad e \lor e \leftarrow d \}. \]

For any nonempty set \( S \subseteq \{a_1, \ldots, a_n\} \), \( S \cup \{b, d\} \) is a loop of \( \text{remove}(\Pi_2, a) \). Then \( \text{Loop}(\text{remove}(\Pi_2, a)) = 4 + n + 2^n \). \( \text{HC}(\text{remove}(\Pi_2, a)) = \{a, d, a_1, \ldots, a_n\} \).

**Lemma 3** Let \( \Pi \) be a DLP and \( E \subseteq \text{Atoms}(\Pi) \). Then for each \( L \in \text{Loop}(\text{remove}(\Pi, E)) \), either \( L \cap E = \emptyset \) or \( L = \{a\} \) for some \( a \in E \).

From Theorem 1 and above lemmas, we get our main result.

**Theorem 4** For any DLP \( \Pi \),

\[ A^S(\Pi) = A^S(\text{sh}(\text{remove}(\Pi, \text{HC}(\Pi))). \]

Then we can equivalently rewrite a DLP \( \Pi \) to the NLP \( \text{sh}(\text{remove}(\Pi, \text{HC}(\Pi))). \)

Now we discuss some optimizations for the new rewriting. Given a DLP \( \Pi \) with \( n \) different atoms and \( m \) different rules, (the number of rules) \( |\text{remove}(\Pi, a)| \) for an atom \( a \) is \( m + m^2 \) in the worst case, \( |\text{remove}(\Pi, \text{HC}(\Pi))| = O(n^2m) \), and \( |\text{sh}(\text{remove}(\Pi, \text{HC}(\Pi)))| = O(n \cdot m^2) \). As proved in [Eiter et al., 2004], any answer-set-preserving rewriting of a DLP to an NLP must lead in general to an exponential blow-up, providing the polynomial hierarchy does not collapse. However, when \( \Pi \) is “almost” HCF, the rewriting could be efficient. In particular, if \( |\text{HC}(\Pi)| = k \) for a constant \( k \), then \( |\text{sh}(\text{remove}(\Pi, \text{HC}(\Pi)))| = O(n \cdot m^{2k}) \), which is polynomial for the size of \( \Pi \), i.e., \( n \) and \( m \).

We can further reduce the number of unfolded atoms by considering only special loops.

Firstly, \( \text{HC}(\Pi) \) can be replaced by one of its subsets in the rewriting. For a DLP \( \Pi \), we denote

\[ \text{HC}^*(\Pi) = \{ a \mid \text{there exist } r \in \Pi \text{ and } L \in \text{Loop}(\Pi) \text{ such that } |\text{head}(r) \cap L| > 1, r \in R^*(L, \Pi), \quad \text{and } a \in \text{head}(r) \cap L \}. \]
Clearly, $HC^*(\Pi) \subseteq HC(\Pi)$. From Theorem 1, if $HC^*(\Pi) = \emptyset$, then $AS(\Pi) = AS(sh(\Pi))$, even when $\Pi$ is not HCF. So the class of HCF programs can be extended to HCF$^*$ programs with $HC^*(\Pi) = \emptyset$, for which shifting is answer-set-preserving. We have the following corollary of Theorem 4.

**Corollary 5** For any DLP $\Pi$,

$$AS(\Pi) = AS(sh(remove(\Pi, HC^*(\Pi))))$$

Secondly, based on the notions of elementary loops and proper loops, $HC^*(\Pi)$ can be further reduced to some of its subsets in the rewriting.

Gebser et al. [2011] showed that not all loops are necessary for identifying the answer sets among the models of a DLP. They introduced the subclass elementary loops of loops, and refined Theorem 1 by considering elementary loops only. Later, Ji et al. [2015] introduced the subclass proper loops of elementary loops and refined Theorem 1 by considering a special form of loop formulas for proper loops only. As recognizing elementary loops and proper loops for DLPs are coNP-complete, Ji et al. [2015] introduced weak elementary loops and weak proper loops which can be identified in polynomial time.

Based on these notions, we can further refine $HC^*(\Pi)$ in the rewriting. For a DLP $\Pi$, we denote

$$HEL(\Pi) = \{a \mid \text{there exist } r \in \Pi \text{ and an elementary loop } L \text{ of } \Pi \text{ such that } |head(r) \cap L| > 1, r \in R^-(L, \Pi), \text{ and } a \in head(r) \cap L\}.$$  

$HPL(\Pi)$, $HWEL(\Pi)$, and $HWPL(\Pi)$ denote the sets specified by the definition of HEL(\Pi) by replacing “elementary loop” with “proper loop”, “weak elementary loop”, and “weak proper loop”, respectively. Then we have the following corollary of Theorem 4.

**Corollary 6** For any DLP $\Pi$,

$$AS(\Pi) = AS(sh(remove(\Pi, HEL(\Pi))))$$

$$= AS(sh(remove(\Pi, HPL(\Pi))))$$

$$= AS(sh(remove(\Pi, HWEL(\Pi))))$$

$$= AS(sh(remove(\Pi, HWPL(\Pi))))$$

**4 Reducing Loop Formulas for Computing Answer Sets of Rewritten Programs**

From Theorem 1, answer sets can be identified from models of the program that satisfy loop formulas of all loops. In this section, we show that we only need to consider a subset of loops of an unfolded program, i.e., remove($\Pi, E$) for a DLP $\Pi$ and $E \subseteq Atoms(\Pi)$, for obtaining answer sets. This result could be beneficial for SAT-based answer set programming solvers such as ASSAT [Lin and Zhao, 2004], cmodels [Giunchiglia et al., 2006], and clasp [Gebser et al., 2007], for computing answer sets of remove($\Pi, E$) and sh(remove($\Pi, E$)).

We first provide a lemma to specify relations of loops and loop formulas between a DLP $\Pi$ and the unfolded program remove($\Pi, a$) for an atom $a$.

**Lemma 4** Let $\Pi$ be a DLP and $a$ an atom in $\Pi$. Then the following statements hold:

- $\{L \mid L \in Loop(\Pi), a \notin L\} \subseteq Loop(remove(\Pi, a))$.
- $\{L \setminus \{a\} \mid L \in Loop(\Pi) \setminus \emptyset \subseteq Loop(remove(\Pi, a))\}$.
- For any set $L \subseteq Atoms(\Pi)$, if $a \in L$ then $LF(L, \Pi) \equiv LF(L, remove(\Pi, a))$.
- For any set $L \subseteq Atoms(\Pi)$, if $a \notin L$ then
  - $tr(\Pi) \cup (LF(L, remove(\Pi, a)) \cup LF(L, \Pi))$
  - $(LF(L \cup \{a\}, \Pi) \cup LF(L, \Pi))$

Given a DLP $\Pi$ and $E \subseteq Atoms(\Pi)$, we denote

$$Loop^*(\Pi, E) = \bigcup_{a \in E} \{a\} \cup \{L \mid L \cap E = \emptyset \text{ and there exists } S \subseteq E \text{ such that } L \cup S \in Loop(\Pi)\}.$$  

From Lemma 4, $Loop^*(\Pi, E) \subseteq Loop(remove(\Pi, E))$. remove($\Pi, E$) may have loops that are not in $Loop^*(\Pi, E)$. For some DLPs, the size of $Loop(remove(\Pi, E))$ can be exponentially larger than the size of $Loop^*(\Pi, E)$.

**Example 2 (Continued)** $Loop^*(\Pi_2, \{a\}) = Loop(\Pi_2) \setminus \{\{a, b\}\}$ and $|Loop(remove(\Pi_2, a)) \setminus Loop^*(\Pi_2, \{a\})| = 2^n - 1$.

The following theorem shows that, we only need to consider loops in $Loop^*(\Pi, E)$ for identifying answer sets among the models of $tr(remove(\Pi, E))$.

**Theorem 7** Let $\Pi$ be a DLP and $E \subseteq Atoms(\Pi)$. Then $S \in AS(remove(\Pi, E))$ iff $S \cup \{-p \mid p \notin S\}$ is a model of

$$tr(\Pi) \land \bigwedge_{L \in Loop^*(\Pi, E)} LF(L, remove(\Pi, E)).$$

Recalling Theorem 1, we can improve the computation of answer sets for remove($\Pi, E$) on SAT-based answer set programming solvers by only considering loop formulas for loops in a subset of $Loop(remove(\Pi, E))$.

**Corollary 8** Let $\Pi$ be a DLP and $E \subseteq Atoms(\Pi)$. Then $S \in AS(sh(remove(\Pi, E)))$ iff $S \cup \{-p \mid p \notin S\}$ is a model of

$$tr(\Pi) \land \bigwedge_{L \in Loop^*(\Pi, E)} LF(L, sh(remove(\Pi, E))).$$

Theorem 7 and Corollary 8 show that although our transformation to NLP may result in an exponential blow-up, there is not necessarily a blow-up in loop formulas for the computation. In particular, $|Loop^*(\Pi, E)| \leq |Loop(\Pi)|$ for any set $E \subseteq Atoms(\Pi)$, then the number of loops that need to be considered for computing answer sets of remove($\Pi, E$) or sh(remove($\Pi, E$)) is not larger than that of $\Pi$.

**5 Approximating a DLP by Restricted Unfolding**

In this section, we consider another application of the new rewriting for approximating a DLP $\Pi$ by an NLP. First, we specify the notion of approximation.
Definition 3 (Sound Approximation) A DLP $\Pi^*$ is a (sound) approximation of a DLP $\Pi$, if $AS(\Pi^*) \subseteq AS(\Pi)$.

The intuition behind the definition is that, when it is hard to compute answer sets of a DLP $\Pi$, we would like to construct an approximation $\Pi^*$ of $\Pi$ so that $\Pi^*$ can compute some answer sets of $\Pi$ easier. The approximation is sound in the sense that every answer set computed from the approximation is an answer set of $\Pi$.

Definition 4 (Closed Sequence of Sound Approximations) A sequence of DLPs $(\Pi_1, \ldots, \Pi_n)$ is a closed sequence of (sound) approximations of a DLP $\Pi$ if $AS(\Pi_i) \subseteq AS(\Pi_{i+1})$ for each $1 \leq i \leq n - 1$, and $AS(\Pi_n) = AS(\Pi)$.

Intuitively, given a closed sequence of sound approximations of a DLP $\Pi$, we can construct an anytime algorithm by computing answer sets of $\Pi_i$’s one by one, which is expected to find better solutions the more time it keeps running. Eventually, the algorithm would return all answer sets of $\Pi$.

In the following, we specify a closed sequence of sound approximations of a DLP by restricted unfolding and shifting.

Lemma 5 For any DLP $\Pi$ and $a \in \text{Atoms}(\Pi)$, $AS(\text{remove}(\text{sh}(\Pi), a)) \subseteq AS(\text{sh}(\text{remove}(\Pi, a)))$.

Theorem 9 Let $\Pi$ be a DLP and $E_1 \subseteq E_2 \subseteq HC(\Pi)$. Then

$$AS(\text{sh}(\Pi)) \subseteq AS(\text{sh}(\text{remove}(\Pi, E_1))) \subseteq AS(\text{sh}(\text{remove}(\Pi, E_2))) \subseteq AS(\Pi).$$

From Theorem 9, we can construct a closed sequence of sound approximations of a DLP.

Corollary 10 Let $\Pi$ be a DLP. $HC(\Pi) = \{a_1, \ldots, a_n\}$. Then

$$(\text{sh}(\Pi), \text{sh}(\text{remove}(\Pi, \{a_1\})), \text{sh}(\text{remove}(\Pi, \{a_1, a_2\})), \ldots, \text{sh}(\text{remove}(\Pi, \{a_1, \ldots, a_n\})))$$

is a closed sequence of sound approximations of $\Pi$.

Example 1 (Continued) $AS(\text{sh}(\Pi_1)) = \emptyset$ and $AS(\Pi_1) = \text{AS}((\text{rem}(\Pi_1, a)) = AS(\text{sh}(\text{remove}(\Pi_1, a))) = \{(a, b, c, d, e)\}$.

Now we provide an anytime algorithm (Algorithm 1) to compute answer sets of a DLP by using NLP solvers.

Algorithm 1: Computing answer sets of a DLP $\Pi$

1. $H := HC(\Pi)$;
2. $as := AS(\text{sh}(\Pi))$;
3. output $as$; \(//\) current computed answer sets of $\Pi$
4. for each $a \in H$
5. \quad $as := \text{remove}(\Pi, a)$;
6. \quad $as := AS(\text{sh}(\Pi))$;
7. output $as$;
8. return $as$;

Theorem 11 For any DLP $\Pi$, Algorithm 1 outputs a subset of $AS(\Pi)$ at each step and returns $AS(\Pi)$ at last.

We can apply the approaches discussed in previous sections to optimize Algorithm 1 as well. In particular, we can replace $HC(\Pi)$ by $HC^*(\Pi)$ or $HVVEL(\Pi)$ in the algorithm. We can also use the approach proposed in Section 4 to improve the computation of answer sets for $\text{sh}(\Pi)$.

Algorithm 1 provides an approach to compute an answer set of a DLP by calling NLP solvers. Janhunen et al. [2006] also provided such an approach. In their approach, a DLP $\Pi$ is first rewritten to an NLP $\text{Gen}(\Pi)$ to produce candidate models $S$, then another NLP $\text{Test}(\Pi, S)$ is constructed to check whether $S$ is an answer set of $\Pi$. The rewriting proposed in [Janhunen et al., 2006] is different from our rewriting in the sense that $AS(\Pi) \subseteq AS(\text{Gen}(\Pi))$, i.e., $\text{Gen}(\Pi)$ is not a sound approximation of $\Pi$, while we provide a closed sequence of sound approximations.

6 Some Experiments

We have implemented a program\(^1\) for rewriting a DLP $P$ to an NLP $P^n$, where $P^n$ is a simplified program from $\text{sh}(\text{remove}(P, HC(P)))$ by removing redundancies like rules $r$ with $\text{head}(r) \cap \text{body}^+(r) \neq \emptyset$ or $\text{body}^+(r) \cap \text{body}^-(r) \neq \emptyset$. By Theorem 4, $AS(P) = AS(P^n)$. Intuitively, if the size of the NLP $P^n$ is similar to the size of the DLP $P$, it would be more efficient to compute answer sets of $P^n$ using NLP solvers than it is to compute answer sets of $P$ using DLP solvers. To verify this, we tried our program on a number of benchmarks.

First, we tried benchmarks that were frequently used to compare the performance of DLP solvers [Denecker et al., 2009; Gebser et al., 2013]. However, for most non-HCF DLPs $P$ in these benchmarks, the size of $P^n$ is much larger than the size of $P$. Then it is harder to compute answer sets of these $P^n$ using NLP solvers.

Next, we tried some programs constructed from Niemel"a’s [1999] encoding of $H$ of the Hamiltonian Circuit (HC) problem. Specifically, we construct a program $Q_0$:

$$a \lor b \leftarrow \text{reached}(1). \quad a \leftarrow b. \quad b \leftarrow a.$$  

where $\text{reached}(1)$ is an atom in $H$ and $a, b$ are new atoms not appearing in $H$. We consider programs $Q = Q_0 \cup H \cup G$ for corresponding graphs $G$ of the HC problem in the experiment. Let $Q^n_0$ be the program:

$$b \leftarrow \text{reached}(1), \text{not } a. \quad a \leftarrow \text{reached}(1), \text{not } b.$$

We denote $Q^n = Q^n_0 \cup H \cup G$. Clearly, $Q^n = \text{sh}(\text{remove}(Q, HC(Q)))$.

In Table 1, we compare the running times of computing an answer set of $Q$ and $Q^n$ by corresponding solvers\(^2\). In the table, $\text{rand}_m(t)$ stands for a class of randomly generated $t$ graphs with $m$ vertices and $k$ arcs. The numbers under “$Q$” refer to the average running times for computing an answer set of corresponding programs $Q$, using corresponding

\(^1\)Our implementation is available on the web http://ss.sysu.edu.cn/~weh/dlp2nlp.html.

\(^2\)Our experiments were done on an Intel(R) Core(TM) i7-2600 3.40GHz CPU and 32GB RAM. The reported times are in CPU seconds as reported by Linux “/usr/bin/time” command.
DLP solvers, i.e., clasp\(^3\) (version 3.1.4) and cmodels\(^4\) (version 3.8.5), and those under “\(Q^n\)” refer to the average running times for computing an answer set of corresponding programs \(Q^n\), using corresponding NLP solvers, i.e., clasp and cmodels. As can be seen, for these programs, computing an answer set of \(Q^n\) by NLP solvers is more efficient.

<table>
<thead>
<tr>
<th>Graphs</th>
<th>Q</th>
<th>Q(^n)</th>
</tr>
</thead>
<tbody>
<tr>
<td>rand\textunderscore 200,1800(5)</td>
<td>1.47</td>
<td>5.14</td>
</tr>
<tr>
<td>rand\textunderscore 250,2200(5)</td>
<td>43.91</td>
<td>6.67</td>
</tr>
<tr>
<td>rand\textunderscore 300,2700(5)</td>
<td>616.76</td>
<td>332.80</td>
</tr>
<tr>
<td>rand\textunderscore 350,3000(5)</td>
<td>1102.61 (3)</td>
<td>1080.56 (3)</td>
</tr>
</tbody>
</table>

When a solver could not compute an answer set of a program in 1800 seconds, we count the running time as 1800 seconds. We also report the number of such programs in the parentheses.

Table 1: Comparing \(Q\) and \(Q^n\)

7 Eliminating Disjunctions for First-order Programs on Finite Structures

In this section, we extend the new rewriting to non-ground programs with first-order answer set semantics on finite structures. We first review some basic notions.

In the following, we consider the first-order disjunctive logic programs [Ferraris et al., 2011], with similar syntactic forms as DLP programs. A distinction is made between predicates occurring only in rule bodies and the reminder predicates. A predicate \(P\) in a first-order DLP \(Π\) is intensional (or IDB) if \(P\) occurs in \(head(r)\) for some \(r \in Π\), and extensional (or EDB) otherwise. We use \(Π(P)\), \(Π_DB(P)\), and \(Π_{EDB}(P)\) to denote sets of predicates, intensional predicates, extensional predicates occurring in \(Π\) respectively. As with [Lee and Meng, 2008], we assume w.l.o.g. that \(Π\) is in normal form without object constants occurring in rule heads.

We use first-order (ground) substitution, unifier, most general unifier (mgu) defined in the standard way. In contrast to grounded DLP, the answer sets of a first-order DLP is defined using first-order structures. In particular, a finite structure \(M\) is a tuple \((D, c^M_1, ..., c^M_m, P^M_1, ..., P^M_n)\), where \(D\) is a finite set called the domain of \(M\), \(c^M_i \in A\) (the interpretation of constant \(c_i\)) \((1 \leq i \leq m)\), and \(P^M_j\) (the interpretation of a \(k\)-ary predicate symbol \(P_j\)) \((1 \leq j \leq n)\) a \(k\)-ary relation on \(D\). We refer the readers to [Asuncion et al., 2012] for the definition of answer sets of first-order DLP on finite structures.

The answer set semantics of a first-order DLP on finite structures can be captured by first-order loop formulas [Chen et al., 2006; Lee and Meng, 2008] as well.

For a first-order DLP \(Π\), the positive dependency graph of \(Π\), written \(G_{PD}\), is the infinite graph \((\Gamma, E)\), where \(\Gamma\) is the set of atoms constructed from \(Π_DB(P)\), \((a, b) \in E\) if there is a rule \(r \in Π\) and a substitution \(θ\) such that \(P(\tilde{t})\ θ = a\) for some \(P(\tilde{t})\ θ = b\) for some \(Q(\tilde{t})\ θ = body^+(r)\). A finite nonempty subset \(L\) of \(Π\) is called a (first-order) loop of \(Π\) if the subgraph of \(G_{PD}\) induced by \(L\) is strongly connected. Specially, for each atom \(a \in V\), \(\{a\}\) is a loop of \(Π\). We use Loop(\(Π\)) to denote the set of loops of \(Π\).

Given a loop \(L\) of a DLP \(Π\) in normal form, we first rename variables in \(Π\) so that no variables of \(Π\) occur in \(L\). Then the (first-order) external support formula of \(L\) for \(Π\), denoted by \(ES(Π, L)\), is the following disjunction

\[
\bigvee_{\theta, r : r \in L, head(r) \theta \not\in L} (\bigwedge \{P(\tilde{t}) \mid P(\tilde{t}) \in L \land \bigwedge \{P(\tilde{t}) \mid P(\tilde{t}) \not\in L \land (\tilde{t} \not\in \tilde{v})\} \bigwedge \neg \bigvee \{P(\tilde{t}) \mid P(\tilde{t}) \not\in L \land (\tilde{t} \not\in \tilde{v})\})
\]

The (first-order) loop formula of \(L\) in \(Π\), denoted by \(LF(Π, L)\), is the universal closure of \(\bigwedge_{\theta = \theta_1 \cup \theta_2 : \theta_1 \not\subseteq \theta_2} ES(Π, L)\).

We use \(LF(Π)\) to denote the set of all loop formulas in \(Π\).

Theorem 12 (Proposition 4 in [Lee and Meng, 2008]) Let \(Π\) be a DLP and \(M\) a finite structure. \(M\) is an answer set of \(Π\) if and only if \(M\) is a model of \(\{LF(Π)\} \cup \{LF(Π)\}\).

Given a first-order DLP \(Π\), the shifting result \(sh(Π)\) is defined in the same way as in the propositional case. Every answer set of \(sh(Π)\) is also an answer set of \(Π\), but the converse is not true in general. Similarly, a first-order DLP \(Π\) is called “Head-Cycle-Free” (HCF), if there is no loop \(L\), rule \(r\) in \(Π\), and substitution \(θ\) such that \(|L \cap head(r)\ θ| > 1\).

Based on the progression semantics of answer sets [Zhou and Zhang, 2011], Zhou [2015] provided the following proposition.

Proposition 1 Let \(P\) be an HCF first-order DLP. Then, \(ASF(sh(Π)) = AS(Π)\).

Definition 5 (Elementary unfolding in first-order case)

Given two rules \(r_1\) and \(r_2\) (with suitable variable renaming), there are atoms \(a \in head(r_1)\) and \(b \in body^+(r_2)\) such that \(a\) and \(b\) are unifiable with an mgu \(θ\). The rule obtained by elementary unfolding \(r_2\) at \(r_1\) at \(a\), denoted by unfold\((r_1, r_2, a, b)\), is the rule \(r\) such that \(head(r) = ((head(r_1) \setminus \{a\}) \cup head(r_2))\ θ\) and \(body(r) = (body(r_1) \cup body(r_2) \setminus \{b\})\ θ\).

Let \(Π\) be a DLP, \(r\) a rule in \(Π\), and \(a\) an atom in \(body^+(r)\). We denote \(U_{unf}(Π, r, b)\) to be the set of rules constructed by the following procedure:

\[
S_0 := \{(r', a') \mid a' \in head(r')\text{ such that }a'\text{ and }b\text{ are unifiable}\};
\]

\(i := 0\);

\(S_{i+1} := \emptyset\);

\(R_{i+1} := \emptyset\);

for each \((r', a') \in S_i\), if \(a'\) and \(b\) are unifiable

\[
r'' := \text{unfold}(r', r, a', b);
\]

\(R_{i+1} := R_{i+1} \cup \{r''\}\);

\(θ := \text{the mgu that }a'\text{ and }b\text{ are unifiable with;}
\]

\(S_{i+1} := S_i \cup \{(r'', a'' \ θ) \mid (r', a') \in S_i, a'' \neq a'\}\);

\(i := i + 1\);

\(\)
Lemma 6\head

\textit{The following procedure:} due to the unfolding process. Zhou [2015] identified a class having the same sets of answer sets. However, they are equivalent in the sense of choice-bounded DLPs. Let $\Pi$ be a DLP and $a$ an atom in head$(r)$ for some $r \in \Pi$. We denote remove$(\Pi, a)$ to be a set of rules constructed by the following procedure:

\begin{itemize}
  \item $R := \Pi$;
  \item while there is a rule $r \in R$ and $b \in \text{body}^+(r)$ such that $a$ and $b$ are unifiable
    \begin{itemize}
      \item $R := \text{remove}(R, r, b)$;
    \end{itemize}
  \end{itemize}

remove$(\Pi, a) := R$.

Note that, by choosing different rules and corresponding atoms in the procedure, remove$(\Pi, a)$ would return different programs. However, they are equivalent in the sense of having the same sets of answer sets.

The above procedure may not terminate for some DLPs due to the unfolding process. Zhou [2015] identified a class of first-order DLPs called choice-bounded and showed that the process of applying the unfolding transformation when it is possible would always terminate for choice-bounded DLPs. Then for a choice-bounded DLP $\Pi$, the above procedure terminates and remove$(\Pi, a)$ always exists.

Theorem 13\head

Let $\Pi$ be a DLP, $a$ an atom in head$(r)$ for some $r \in \Pi$, and remove$(\Pi, a)$ exists. Then $AS(\Pi) = AS(\text{remove}(\Pi, a))$.

Given a DLP $\Pi$, we denote head$(\Pi) = \{a \mid a \in \text{head}(r) \text{ for some } r \in \Pi\}$. Let $\Pi$ be a DLP and a set $A \subseteq \text{head}(\Pi)$. We denote remove$(\Pi, A)$ to be a set of rules constructed by the following procedure:

\begin{itemize}
  \item $R := \Pi$;
  \item for each $a \in A$
    \begin{itemize}
      \item $R := \text{remove}(R, a)$;
    \end{itemize}
  \end{itemize}

remove$(\Pi, A) := R$.

Corollary 14\head

Let $P$ be a DLP, $A \subseteq \text{head}(\Pi)$, and remove$(\Pi, A)$ exists. Then $AS(\Pi) = AS(\text{remove}(\Pi, A))$.

Now we extend our rewriting to first-order DLPs. For a DLP $\Pi$, we denote

$HC(\Pi) = \{a \mid a \in \text{head}(r), \text{ there exists } r \in \Pi, \text{ a loop } L \text{ of } \Pi, \text{ and a substitution } \theta \text{ such that } |L \cap \text{head}(r)\theta| > 1, a \in \text{head}(r) \text{ and } a\theta \in L\}$.

We provide the following lemma to draw a sketch of the proof for the main result.

Lemma 6\head

Let $\Pi$ be a DLP, $A \subseteq \text{head}(\Pi)$, and remove$(\Pi, A)$ exists. Then the following statements hold:

- A finite structure $M$ satisfies $LF(\Pi)$ and $tr(\Pi)$ if and only if $M$ satisfies $LF(\text{remove}(\Pi, A))$ and $tr(\text{remove}(\Pi, A))$, where $LF(\Pi) = \{LF\{a\}', \Pi \mid a' = a\theta \text{ for some } a \in A \text{ and substitution } \theta\}$.
- For each $L \in \text{Loop}(\text{remove}(\Pi, A))$, if there does not exist a substitution $\theta$ such that $L \cap HC(\Pi)\theta \neq \emptyset$, then $LF(L, \text{remove}(\Pi, A)) \equiv LF(L, \text{sh}(\text{remove}(\Pi, A)))$.
- for each $L \in \text{Loop}(\text{remove}(\Pi, A))$, either there does not exist a substitution $\theta$ with $L \cap A\theta \neq \emptyset$ or $L = \{a\}$ for some $a' \in A$ and substitution $\sigma$ with $a'\sigma = a$.

Theorem 15\head

Let $\Pi$ be a DLP and remove$(\Pi, HC(\Pi))$ exists. Then $AS(\Pi) = AS(\text{sh}(\text{remove}(\Pi, HC(\Pi))))$.

Asuncion et al. [2012] provided a solver of first-order NLPs by translating them to first-order sentences on finite structures. From Theorem 15, we can use the solver to compute answer sets of some first-order DLPs.

8 Conclusion\head

In this paper, we provide a new answer-set-preserving rewriting of a DLP to an NLP by first apply the unfolding transformation on “culprit” atoms, i.e., atoms that prevent the program from being HCF, then shift the resulting program. This result shows that the complexity penalty for extending NLPs by including disjunctive rules is closely tied to these “culprit” atoms. Hence, different from other transformations for eliminating disjunctions in literature, our rewriting is efficient for “almost” HCF programs, i.e., DLPs that have only a few “culprit” atoms. Later, we show that the computation of answer sets for a rewritten program can be improved by reducing the number of loops that need to be considered and the size of these loops is not larger than that of the original program. This result shows that although our transformation to NLP may result in an exponential blow-up, there is not necessarily a blow-up in loop formulas for the computation. Based on the new writing, we provide an anytime algorithm to compute answer sets of a DLP by calling NLP solvers. At last, we extend the rewriting to non-ground answer set programs on finite structures.

9 Acknowledgments\head

We thank the reviewers for their comments and suggestions for improving the paper. We thank Xiaoping Chen and his research group for their useful discussions. Jianmin Ji’s research was partially supported by the National Natural Science Foundation of China under grant 61175057, the National Natural Science Foundation for the Youth of China under grant 61403359, as well as the USTC Key Direction Project and the USTC 985 Project. Hai Wan’s research was in part supported by the National Natural Science Foundation of China under grant 61573386, Natural Science Foundation of Guangdong Province under grant 508257922126, Guangdong Province Science and Technology Plan projects under grant 509208496103, and Sun Yat-sen University Young Teachers Cultivation Project under grant 16lgyy40. This work is also partially supported by ARC under grant DP130102302.
References


